

The intermediate logics on the second slice

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In [1], Hosoi defined the notion *slice* and proved that for any intermediate logic M there exists a unique slice \mathcal{S}_n ($n \leq \omega$) such that $M \in \mathcal{S}_n$. In this paper, we will give a complete description of the intermediate logics on \mathcal{S}_2 and those logics on \mathcal{S}_3 that contain the axiom $\neg p \vee \neg \neg p$. We assume a knowledge of the results and the notations in [1] and [2].

For a *model* M , M means not only a logic as a set of formulas but a pseudo-Boolean algebra¹⁾. In [3], McKay proved that any intermediate logic can be represented as an intersection of some intermediate logics of the form $S_1 \uparrow M$ or S_1 . Using this result, we prove, in §1, that if $M \in \mathcal{S}_2$ then $M \supset \subset S_1 \uparrow S_1^\sigma$ for some n ($1 \leq n \leq \omega$). Moreover we will give an axiomatization of each logic in \mathcal{S}_2 . The results in §2 resemble those in §1.

§1.

We first cite a theorem due to McKay [3] without proof.

THEOREM 1.1 (McKay). *If M is an intermediate logic, then there are models N_λ ($\lambda \in A$) such that $M \supset \subset \bigcap_{\lambda \in A} (S_1 \uparrow N_\lambda)$.*

Remark that, if M is a finite model, then A and all N_λ 's are finite. In general, the cardinal of N_λ is equal or less than that of M .

Now, using the above theorem we investigate the intermediate logics on the second slice \mathcal{S}_2 . By 1.1, if $M \in \mathcal{S}_2$, then $M \supset \subset \bigcap_{\lambda \in A} (S_1 \uparrow N_\lambda)$ for some N_λ 's. Since $M \subset S_1 \uparrow N_\lambda$ for any $\lambda \in A$, $S_1 \uparrow N_\lambda \in \mathcal{S}_m$ for some $m \leq 2$ and therefore $S_1 \uparrow N_\lambda \in \mathcal{S}_2$ and $N_\lambda \in \mathcal{S}_1$ by Theorem 6.2 in [1]. This means that each model N_λ is a Boolean algebra. Notice that each S_1^σ is isomorphic to the field of all subsets of $\{\mu \mid \mu < \sigma\}$, where σ is any cardinal number. $S_1 \uparrow S_1^\sigma$ is an example of a model in \mathcal{S}_2 . We write M_σ for $S_1 \uparrow S_1^\sigma$ for a cardinal number σ . It is easy to see that if $m < \omega$ then the cardinal of M_m is $2^m + 1$.

Now, we have

LEMMA 1.2. *Let M be a Boolean algebra whose cardinal is σ . Then there exists a subalgebra N of M such that N is isomorphic to S_1^n ($n < \omega$), if $2^n \leq \sigma$.*

COROLLARY 1.3. *If $m < \omega$ and $m < n$, then $M_n \not\subseteq M_m$. (n may be infinite.)*

¹⁾ That is, $A \in M$ means that a formula A is valid in M . But $a \in M$ means that a is an element of a pseudo-Boolean algebra M . As for pseudo-Boolean algebras, see e.g. [5].

PROOF. By 1.2, there is a subalgebra N of S_1^σ such that N is isomorphic to S_1^m . So M_m is isomorphic to a subalgebra of M_n . Thus $M_n \subset M_m$. Now let X_n be the formula $\bigvee_{1 \leq i < j \leq n+1} (q_i \equiv q_j)^2$. Then by Corollary 6.10 in [1], $X_{2^{m+1}} \in M_m$ but $X_{2^{m+1}} \notin M_n$. For $r(M_n) \geq r(M_{m+1}) = 2^{m+1} + 1 > 2^m + 1$. Thus $M_m \not\supset M_n$.

LEMMA 1.4. *If σ is an infinite cardinal number, then $\bigcap_{m < \omega} M_m \subset M_\sigma$.*

PROOF. Suppose that a formula A is not in M_σ . Then, there is an assignment f of M_σ such that $f(A) \neq 1_{M_\sigma}$ ³⁾. Let $\{p_1, \dots, p_n\}$ be the set of all propositional variables appearing in A . We write Q for the set $\{f(p_i) \mid f(p_i) \neq 1_{M_\sigma}, 1 \leq i \leq n\}$. Let Q^* be the subalgebra of S_1^σ generated by Q . Since Q is finite, Q^* is also finite. So Q^* is isomorphic to S_1^m for some $m < \omega$. Thus f can be considered as an assignment of $S_1 \uparrow S_1^m$, or M_m . But $f(A) \neq 1_{M_\sigma} = 1_{M_m}$. So, A is not in M_m for some $m < \omega$.

COROLLARY 1.5. *Let M be a countably infinite Boolean algebra. Then $S_1 \uparrow M \supset M_\omega$.*

PROOF. We can show that $M_\omega \supset \bigcap_{m < \omega} M_m$ by 1.2 and 1.4. So $M_\omega \supset \bigcap_{m < \omega} M_m \supset S_1 \uparrow M$ by 1.2. By the Stone representation theorem of Boolean algebra, M is isomorphic to a subalgebra of S_1^σ for some infinite cardinal number σ , since M is countably infinite. So $S_1 \uparrow M \supset S_1 \uparrow S_1^\sigma = M_\sigma \supset \bigcap_{m < \omega} M_m \supset M_\omega$. Therefore $S_1 \uparrow M \supset M_\omega$.

THEOREM 1.6. *If $M \in \mathcal{S}_2$, then $M \supset M_n$ for some n ($1 \leq n \leq \omega$). Hence \mathcal{S}_2 is linearly ordered by \subset .*

PROOF. Let M be in \mathcal{S}_2 . We can take a countable model M_0 such that $M \supset M_0$. By 1.1 and its remark, $M_0 \supset \bigcap_{\lambda \in A} (S_1 \uparrow N_\lambda)$ for some N_λ 's in \mathcal{S}_1 such that each N_λ is countable. If N_λ is finite, then $S_1 \uparrow N_\lambda = M_m$ for some $m < \omega$. Otherwise, $S_1 \uparrow N_\lambda \supset M_\omega$ by 1.5. So, $M \supset \bigcap_{\lambda \in A} M_{m_\lambda}$, where $1 \leq m_\lambda \leq \omega$. Let n be $\sup \{m_\lambda \mid \lambda \in A\}$. Then $1 \leq n \leq \omega$ and $\bigcap_{\lambda \in A} M_{m_\lambda} \supset M_n$. Hence $M \supset M_n$.

Next we will give an axiomatization of each logic in \mathcal{S}_2 . Let T_k ($1 \leq k \leq \omega$) be an intermediate logic obtained by adding two axiom schemata P_2 and X_k to the intuitionistic propositional logic L . (We write $T_k = L + P_2 + X_k$, following the notation in [1].) Let f be a function on natural numbers such that $f(k) = 2^k + 1$.

THEOREM 1.7. *$M_k \supset T_{f(k)}$ if $1 \leq k < \omega$, and $M_\omega \supset LP_2$.*

PROOF. By Corollary 4.7 in [1], LP_2 is the minimal element in \mathcal{S}_2 . So, $M_\omega \supset LP_2$ by 1.6 and 1.3. By the proof of 1.3, $X_{f(k)} \in M_k$. Since $M_k \in \mathcal{S}_2$, $P_2 \in M_k$ and hence $T_{f(k)} \subset M_k$. By 1.6, $T_{f(k)} \supset M_m$ for some m ($k \leq m \leq \omega$). But

²⁾ See [1].

³⁾ 1_M denotes the designated element of a regular model M .

if $m > k$, then $r(M_m) = f(m) > f(k)$ hence $X_{f(k)} \in M_m$. Thus $m = k$. That is, $T_{f(k)} \supset \supset M_k$. We remark that if $f(k) \leq m < f(k+1)$ then $T_m \supset \supset M_k$.

We next show that each logic \mathcal{S}_2 is normalizable (that is, X_k can be interdeducibly expressed by a formula only containing the logical operator \rightarrow)⁴⁾. Define Q_{ij} by $Q_{ij} = (q_i \rightarrow q_j) \rightarrow ((q_j \rightarrow q_i) \rightarrow q_0)$, and Y_n by $Y_n = Q_{12} \rightarrow (Q_{13} \rightarrow \dots \rightarrow (Q_{1, n+1} \rightarrow (Q_{23} \rightarrow (Q_{24} \rightarrow \dots \rightarrow (Q_{n, n+1} \rightarrow q_0) \dots))) \dots)$. Then it is easy to see that X_n and Y_n are interdeducible in L . Hence $M_k \supset \supset L + P_2 + X_{f(k)} \supset \supset L + P_2 + Y_{f(k)}$. Since both P_2 and Y_n contain only the connective \rightarrow , we have the following

COROLLARY 1.8. *Each intermediate logic in \mathcal{S}_2 is normalizable.*

§ 2.

Logics with the axiom $\neg p \vee \neg \neg p$ have been partially studied in [2]. Here we will show models for those logics on \mathcal{S}_3 . (It has been known that only S_1 and S_2 are those in $\mathcal{S}_1 \cup \mathcal{S}_2$ that contain $\neg p \vee \neg \neg p$).

We remark first that $S_1 \uparrow S_1^m \uparrow S_1 \in \mathcal{S}_3$ and $\neg p \vee \neg \neg p \in S_1 \uparrow S_1^m \uparrow S_1$ for any cardinal number m (see 4.8 in [2]).

LEMMA 2.1. *If $\neg p \vee \neg \neg p \in S_1 \uparrow N$ then for any $a, b \in S_1 \uparrow N$ $a \wedge b < \omega$ if $a < \omega$ and $b < \omega$.*

PROOF. Suppose that $a < \omega, b < \omega$ and $a \wedge b = \omega$ for some $a, b \in S_1 \uparrow N$. Let f be an assignment of $S_1 \uparrow N$ such that $f(p) = a$. By the above assumption, $1_N < f(\neg p) \leq b < \omega$. So $1_N < f(\neg \neg p)$ and $1_{S_1 \uparrow N} \neq 1_N \leq f(\neg p \vee \neg \neg p)$. This contradicts $\neg p \vee \neg \neg p \in S_1 \uparrow N$.

We can prove the following lemma similarly as 1.3.

LEMMA 2.2. *If $m < n \leq \omega$ then $S_1 \uparrow S_1^m \uparrow S_1 \not\supset S_1 \uparrow S_1^n \uparrow S_1$.*

LEMMA 2.3. *Let N be a model in \mathcal{S}_2 such that $\neg p \vee \neg \neg p \in S_1 \uparrow N$. Then $\bigcap_{m < \omega} (S_1 \uparrow S_1^m \uparrow S_1) \subset S_1 \uparrow N$. Furthermore, if N is infinite then $\bigcap_{m < \omega} (S_1 \uparrow S_1^m \uparrow S_1) \supset \supset S_1 \uparrow N$.*

PROOF. In this proof we write 2 for the minimal element 1_N of N . At first, we show that the relation

$$\forall a, b (\omega > b > a > 2 \rightarrow a \vee (a \supset b) = 2)$$

holds in N . Assume that there exist a and b such that $\omega > b > a > 2$ and $a \vee (a \supset b) > 2$. We write c for $a \vee (a \supset b)$. Define an assignment g by $g(p_0) = \omega, g(p_1) = b$ and $g(p_2) = c$. Then $f(P_2) = c > 2$ where $P_2 = ((p_2 \rightarrow (((p_1 \rightarrow p_0) \rightarrow p_1) \rightarrow p_1)) \rightarrow p_2) \rightarrow p_2$. This contradicts that $N \in \mathcal{S}_2$. Now, for any fixed b such that $\omega > b > 2$, we can show that the sublattice $\{a | b \geq a \geq 2\}$ of N is a Boolean algebra. Suppose that a

⁴⁾ See [4].

formula $A \in S_1 \uparrow N$. Then there is an assignment f such that $f(A) \neq 1$. Let Q be the finite set $\{f(p_i) | 1 \leq i \leq m\} \cap (N - \{\omega\})$, where p_i 's are all the propositional variables appearing in A . Let $b = \bigwedge_{a \in Q} a$. Since $\neg p \vee \neg \neg p \in S_1 \uparrow N$ and each $a \in Q$ is smaller than ω , we can show that $b < \omega$ by 2.1. Suppose that $b > 2$. We write M for the Boolean algebra $\{a | b \geq a \geq 2\}$ and M^* for the subalgebra of M generated by $Q^{(5)}$. Since Q is finite, so is M^* . Therefore M^* is of the form S_1^m ($m < \omega$). Then it is easy to see that f is also an assignment of $S_1 \uparrow M^* \uparrow S_1$. So, $A \in S_1 \uparrow M^* \uparrow S_1$ and $A \in \bigcap_{m < \omega} (S_1 \uparrow S_1^m \uparrow S_1)$. If $b = 2$ then we can show that f can also be regarded as an assignment of S_3 and hence $A \in S_3$ and $A \in \bigcap_{m < \omega} (S_1 \uparrow S_1^m \uparrow S_1)$. Suppose N is infinite. Then we prove that $S_1 \uparrow S_1^m \uparrow S_1$ is a subalgebra of $S_1 \uparrow N$ for any $m < \omega$. If N is atomic, then there is only one atom by 2.1.⁽⁶⁾ So N is of the form $N_0 \uparrow S_1$ where N_0 is an infinite Boolean algebra. By 1.2, $S_1 \uparrow S_1^m \uparrow S_1$ is a subalgebra of $S_1 \uparrow N_0 \uparrow S_1$. On the other hand, suppose N is not atomic. Now we can take $a (< \omega)$ in N such that there exists an ascending sequence of elements a_i in N . That is, $a < a_1 < a_2 \cdots < \omega$. So, for any $m < \omega$, there is a subalgebra of N isomorphic to $S_1^m \uparrow S_1$ since $\{b | a_i \geq b \geq 2\}$'s are Boolean algebras. Therefore $S_1 \uparrow N \subset S_1 \uparrow S_1^m \uparrow S_1$ for any $m < \omega$. This implies $S_1 \uparrow N \subset \bigcap_{m < \omega} (S_1 \uparrow S_1^m \uparrow S_1)$.

COROLLARY 2.4. $\bigcap_{m < \omega} (S_1 \uparrow S_1^m \uparrow S_1) \supset \subset S_1 \uparrow S_1^{\omega} \uparrow S_1$.

THEOREM 2.5. Let M be a model in \mathcal{S}_3 such that $M \ni \neg p \vee \neg \neg p$. Then there exists $m (\leq \omega)$ such that $M \supset \subset S_1 \uparrow S_1^m \uparrow S_1$.

PROOF. By 1.1, $M \supset \subset \bigcap_{i \in A} (S_1 \uparrow N_i)$. Suppose first that all N_i 's are finite. Then all N_i 's are atomic. As the proof of 2.3, each $S_1 \uparrow N_i$ is of the form $S_1 \uparrow M_i \uparrow S_1$. Since $M_i \in \mathcal{S}_1$ and M_i is finite, M_i is of the form $S_1^{m_i}$ ($m_i < \omega$). Let $m = \sup \{m_i | i \in A\}$. Then by 2.2, $M \supset \subset S_1 \uparrow S_1^m \uparrow S_1$. Suppose next that there is an infinite N_{μ} . Then by 2.3, $\bigcap_{m < \omega} (S_1 \uparrow S_1^m \uparrow S_1) \supset \subset S_1 \uparrow N_{\mu} \supset M$. If $A \in M$ then there is an N_i such that $A \in S_1 \uparrow N_i$. Then there is $m < \omega$ such that $A \in S_1 \uparrow S_1^m \uparrow S_1$ whether N_i is finite or not. So $M \supset \subset \bigcap_{m < \omega} (S_1 \uparrow S_1^m \uparrow S_1)$.

We know that the class $\{M | \neg p \vee \neg \neg p \in M, M \in \mathcal{S}_3\}$ of logics on \mathcal{S}_3 is linearly ordered by \supset , by 2.2 and 2.5.

In [2], LQ_3 is defined as $L + Q + P_3$ where Q is a formula interdeducible with $\neg p \vee \neg \neg p$. Now it is easy to verify

THEOREM 2.6. $LQ_3 \supset \subset S_1 \uparrow S_1^g \uparrow S_1$ and $LQ_3 + X_{g(k)} \supset \subset S_1 \uparrow S_1^k \uparrow S_1$ if $k < \omega$, where $g(k) = 2^k + 2$.

⁵⁾ Remark that if $a < \omega$ then $\neg a = \omega$ by 2.1.

⁶⁾ We say an element a in N is an atom, if $b \geq a$ implies either $b = a$ or $b = \omega$. If for any $b < \omega$ there is an atom a such that $b < a$, we say N is atomic.

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